

Natural Deduction Rules with Iterated Quantifiers

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Iterated Universal Quantifiers

- There are no restrictions on instantiating universal sentences, which means there is no problem with iterated universal quantifiers.
- For example, $(\forall x)(\forall y)Fxy$ can be instantiated as:
 - $(\forall y)Fay$, which in turn can be instantiated as:
 - * Faa , or
 - * Fab , etc.

Iterated Existential Quantifiers

- The instantiation of existential quantifiers is subject to restrictions which rule out some instantiations in the case of multiple quantifiers.
- For example, the following pattern of instantiation is *not* permitted:

1		$(\exists x)(\exists y)Fxy$	P
2		$(\exists y)Fay$	A
3		Faa	A
4			

- The second instantiation violates the restriction that the instantiating name be isolated in the derivation.

Hard Problems

- On page 94, a derivation in 21 steps is given.
- We are told what the two main strategic moves are, but not why they were chosen.
- Reductio is chosen as the basic strategy.
 - The reason is that it would be impossible to use \exists I to get the conclusion.

- This condition usually holds, and so reductio is a good strategy for deriving existential sentences.
- The sub-strategy is to use reductio again to get ‘ $\hat{P}a$ ’ in order to use \forall I.
 - Once again, it is impossible to get this result in any other way.
 - This condition usually holds, so reductio is a good strategy for deriving atomic sentences that cannot be derived using instantiation.